

# Cryptographic Security of ECDSA in Bitcoin



Nicolas T. Courtois

## Dr. Nicolas T. Courtois

1. cryptologist and codebreaker



## UNIVERSITY CIPHER CHAMPION

March 2013



2. payment and smart cards (e.g. bank cards, Oyster cards etc...)



### Oyster cracker vows to clone cards

Cloning kit could sell for just £200, says researcher

Robert Blincoe, vnunet.com, 28 Jul 2008

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 Code Breakers

Members (712)



 IACR Cryptographers





## UCL Bitcoin Seminar

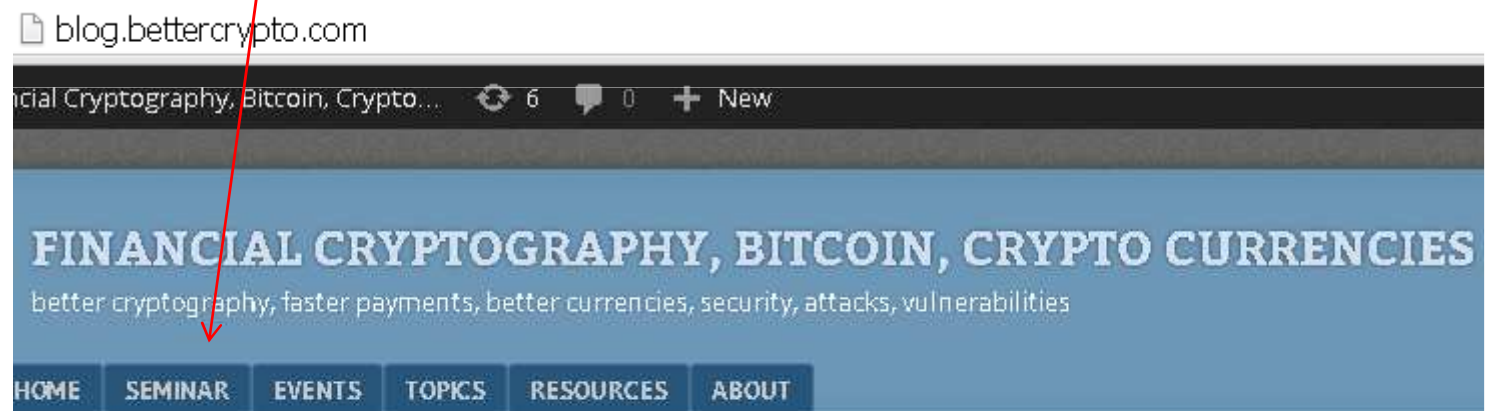
**research** seminar

=>In central London, runs EVERY WEEK!

public web page:

[blog.bettercrypto.com](http://blog.bettercrypto.com) / SEMINAR

or Google "UCL bitcoin seminar"



### New Powerful Attacks On ECDSA In Bitcoin Systems

Posted by admin on 23 October 2014, 10:57 pm

There is a wave of new powerful cryptographic attacks on bitcoin systems.



## My Whole Life:

Tried to improve  
the security baseline...

## My Whole Life:

Tried to improve  
the security baseline...

Crying Wolf!

51%, Elliptic Curve, OpenSSL...



It did NOT help,

The Wolf was allowed to operate



We failed to protect our DATA

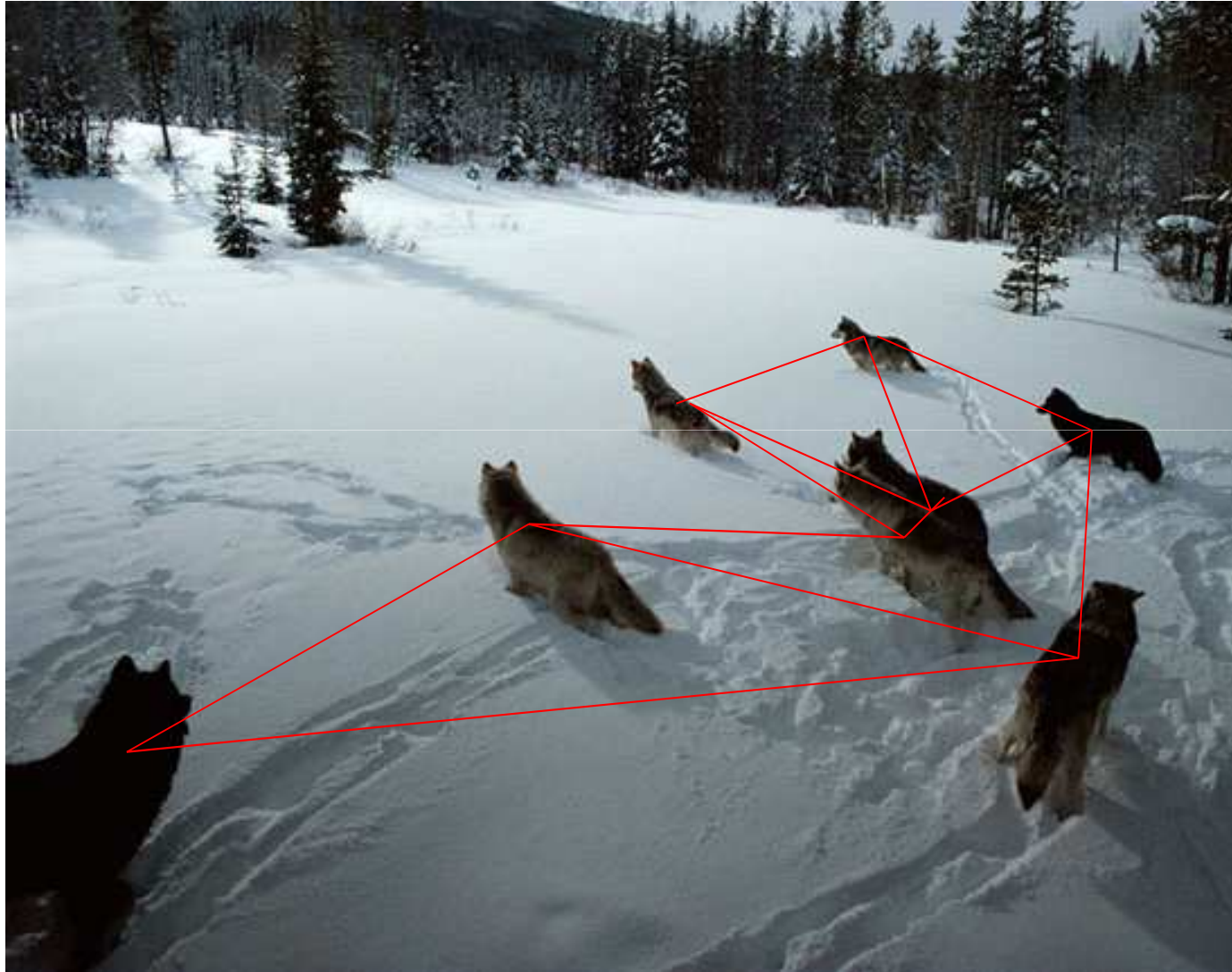




We failed to protect our **MONEY**



# Solution = Decentralized P2P



## Solution = BlockChain



- Until recently, we've needed central bodies – banks, stock markets, governments, police forces – to settle vital questions.
  - Who owns this money?
  - Who controls this company?
  - Who has the right to vote in this election?
- Now we have a small piece of pure, **incorruptible** mathematics enshrined in computer code that will allow people to solve the thorniest problems without reference to “the authorities”.

<http://www.telegraph.co.uk/technology/news/10881213/The-coming-digital-anarchy.html>

[11 June 2014]

**The Telegraph**

The coming digital anarchy

## But Is Cryptography Incorruptible?

NSA 2013 Budget, excerpts:

[...] actively engages the US and foreign IT industries to **covertly influence** and/or overtly leverage their commercial products' designs.

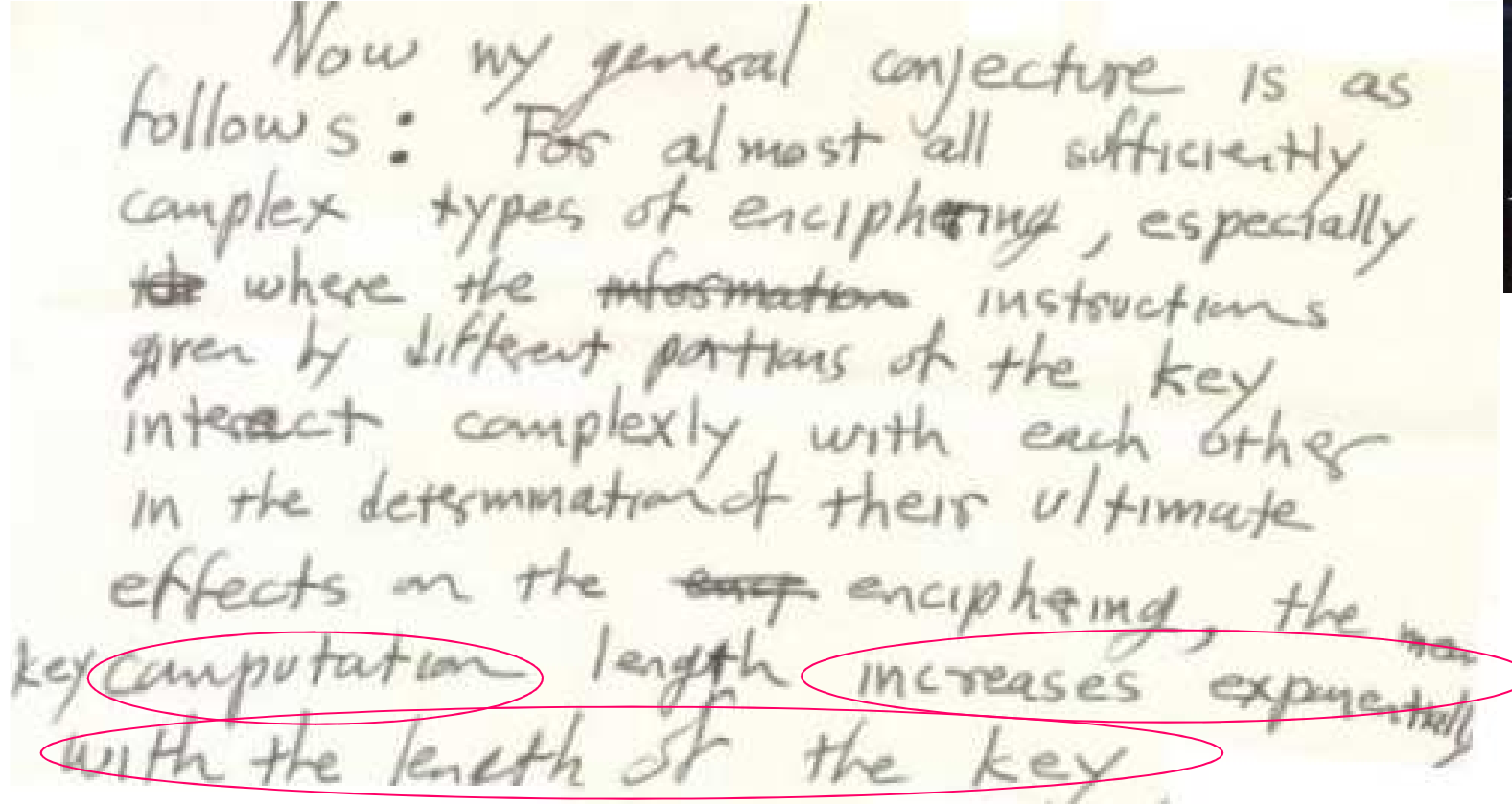


[...] **Insert vulnerabilities** into commercial encryption systems [...]

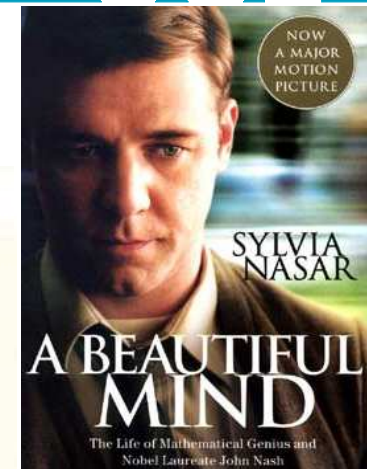
[...] Influence policies, standards and specification for commercial **public key technologies**. [...]

## John Nash - 1955

In 2012 the NSA declassified his hand-written letter:



Now my general conjecture is as follows: For almost all sufficiently complex types of enciphering, especially ~~the~~ where the ~~information~~ instructions given by different portions of the key interact complexly with each other in the determination of their ultimate effects on the ~~enc~~ enciphering, the key computation length increases exponentially with the length of the key.

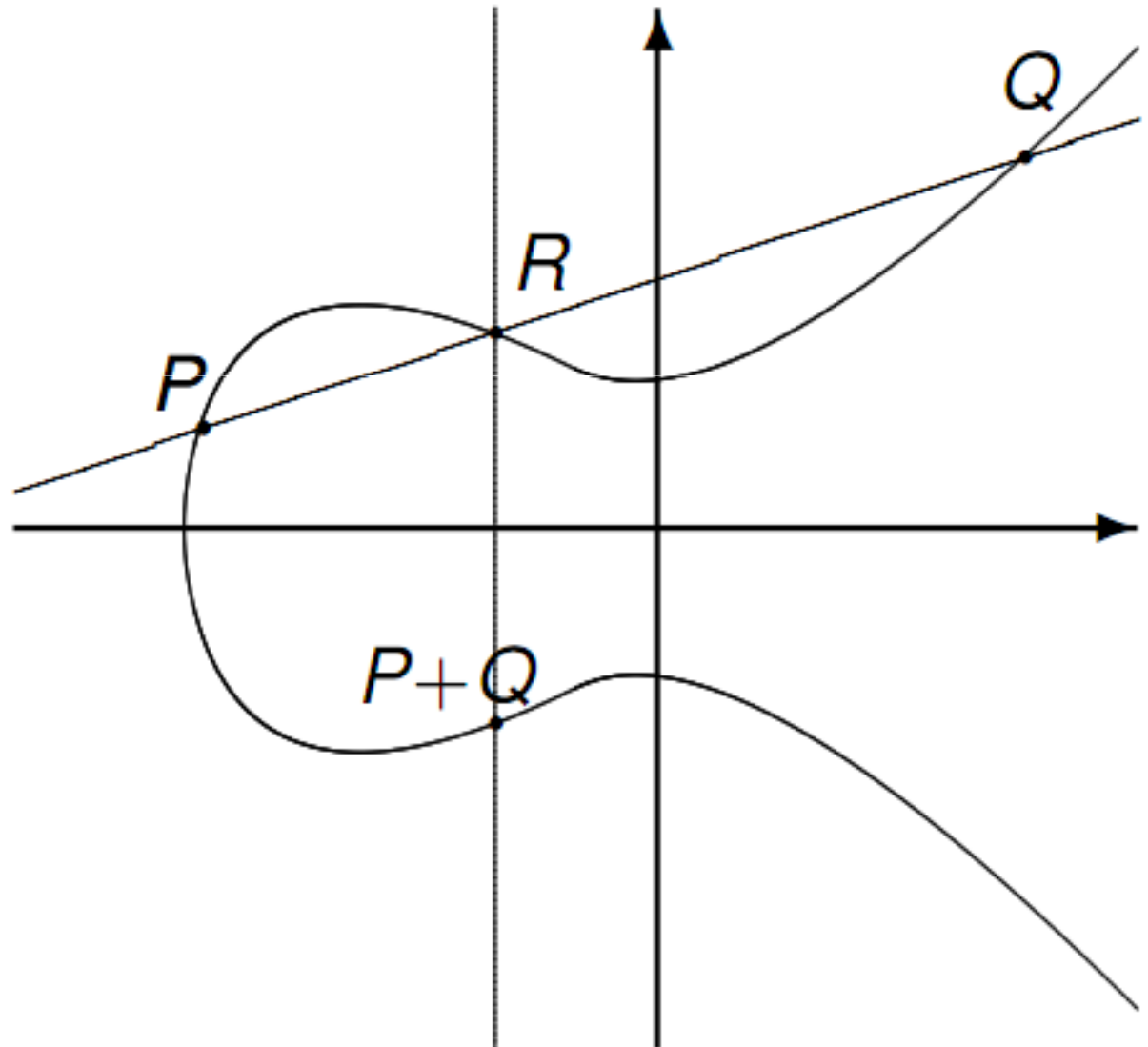
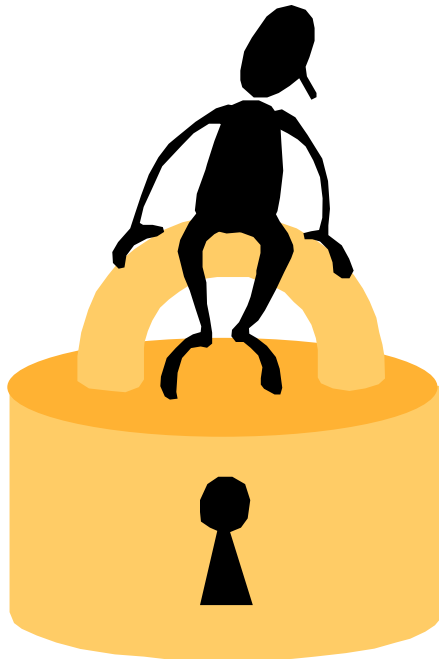


He also says that:

[...] the game of cipher breaking by skilled teams, etc., should become a thing of the past." [...]

# Elliptic Curve Crypto

“exponential  
security”



## ECC - Certicom Challenges [1997, revised 2009]

|          |    |        |          |
|----------|----|--------|----------|
| ECC2K-95 | 97 | 18322  | \$ 5,000 |
| ECC2-97  | 97 | 180448 | \$ 5,000 |

|         |    |       |          |
|---------|----|-------|----------|
| ECCp-97 | 97 | 71982 | \$ 5,000 |
|---------|----|-------|----------|

| Challenge | Field size<br>(in bits) | Estimated number<br>of machine days | Prize<br>(US\$) |
|-----------|-------------------------|-------------------------------------|-----------------|
| ECC2K-108 | 109                     | $1.3 \times 10^6$                   | \$10,000        |
| ECC2-109  | 109                     | $2.1 \times 10^7$                   | \$10,000        |
| ECC2K-130 | 131                     | $2.7 \times 10^9$                   | \$20,000        |
| ECC2-131  | 131                     | $6.6 \times 10^{10}$                | \$20,000        |

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| ECC2-191  | 191                     | $4.07 \times 10^{19}$               | \$40,000        |
| ECC2K-238 | 239                     | $6.83 \times 10^{26}$               | \$50,000        |
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| ECCp-239  | 239                     | $1.4 \times 10^{27}$                | \$50,000        |
| ECCp-359  | 359                     | $3.7 \times 10^{45}$                | \$100,000       |

# TOTAL = 725,000 USD



## P vs. NP

- If you solve P vs. NP it: 1 M\$.
- Nobel price, Abel price in mathematics: roughly 1M\$
- Break bitcoin ECC: About 4 BILLION \$.



# How to Steal Bitcoins

New attacks [Courtois et al. October 2014]

[eprint.iacr.org/2014/848/](http://eprint.iacr.org/2014/848/)

=>more details later...

## Crypto Challenges:

I always liked this idea.

Claiming (very naive) that this would:

“punish those who  
by their ignorance, incompetence  
or because of a hidden agenda,  
put everybody's security at a great risk.”

[Courtois, May 2006, Quo Vadis Cryptology 4 conference]

## ECC - Certicom Challenges [1997, revised 2009]

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**secp256k1**  
**NOT INCLUDED**  
 no price if you  
 break it ☹



## Timely Denial

Dan Brown, chair of SEC [Certicom, Entrust, Fujitsu, Visa International...]

**“I did not know that BitCoin is using secp256k1.  
I am surprised to see anybody use secp256k1 instead of secp256r1”,**

September 2013,

<https://bitcointalk.org/index.php?topic=289795.80>

# Comparison:

| Used/recommended by:                           | secp256k1    | secp256r1            |
|--|--------------|----------------------|
| Bitcoin, anonymous founder, no one to blame... | Y            |                      |
| SEC Certicom Research                          | surprised!   | Y                    |
| TLS, OpenSSL                                   | ever used??? | Y <b>98.3%</b> of EC |
| U.S. ANSI X9.63 for Financial Services         | Y            | Y                    |
| NSA suite B, NATO military crypto              |              | Y                    |
| U.S. NIST                                      |              | Y                    |
| IPSec  |              | Y                    |
| OpenPGP  |              | Y                    |
| Kerberos extension                             |              | Y                    |
| Microsoft implemented it in Vista and Longhorn |              | Y                    |
| EMV bank cards XDA [2013]                      |              | Y                    |
| German BSI federal gov. infosec agency, y=2015 |              | Y                    |
| French national ANSSI agency beyond 2020       |              | Y                    |



## Wanna Bet?

### Bitcoin Cryptography Broken in 2015

Category: [Bitcoin](#)

By  [NCourtois](#) ★★★★★

#### ① Description

The digital signature scheme of bitcoin with SHA256+secp256k1 ECDSA will be broken before 1 September 2015 by cryptography researchers. The attack should allow to forge digital signatures for at least a proportion of 1/1 million bitcoin users and steal money from them. It should be done faster than  $2^{100}$  point additions total including the time to examine the data.



#### ⌚ Decision Logic



bitcoin, cryptography, SHA256, ECDSA, ECDL, secp256k1

<https://www.betmoose.com/bet/bitcoin-cryptography-broken-in-2015-791>

# [betmoose.com](http://betmoose.com) - Totally Anonymous Bets In BTC!

FEATURED

Bitcoin Cryptography Broken in 2015

Category: Bitcoin By NCourtois ★★★★★

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The digital signature scheme of bitcoin with SHA256+secp256k1 ECDSA will be broken before 1 September 2015 by cryptography researchers. The attack should allow to forge digital signatures for at least a proportion of 1/1 million bitcoin users and steal money from them. It should be done faster than  $2^{100}$  point additions total including the time to examine the data.

YES

|            |         |
|------------|---------|
| Volume:    | ₿ 0.140 |
| # of Bets: | 3       |

₿

| PAYOUT | ROI |
|--------|-----|
| ₿ 0.00 | 0%  |

\* assumes current weight and volumes

Place Anonymously

NO

|            |         |
|------------|---------|
| Volume:    | ₿ 0.189 |
| # of Bets: | 6       |

₿ 0.1

| PAYOUT    | ROI    |
|-----------|--------|
| ₿ 0.14327 | 43.27% |

\* assumes current weight and volumes

Place Anonymously



SHA256, ECDSA, ECDL, secp256k1

## Amount?

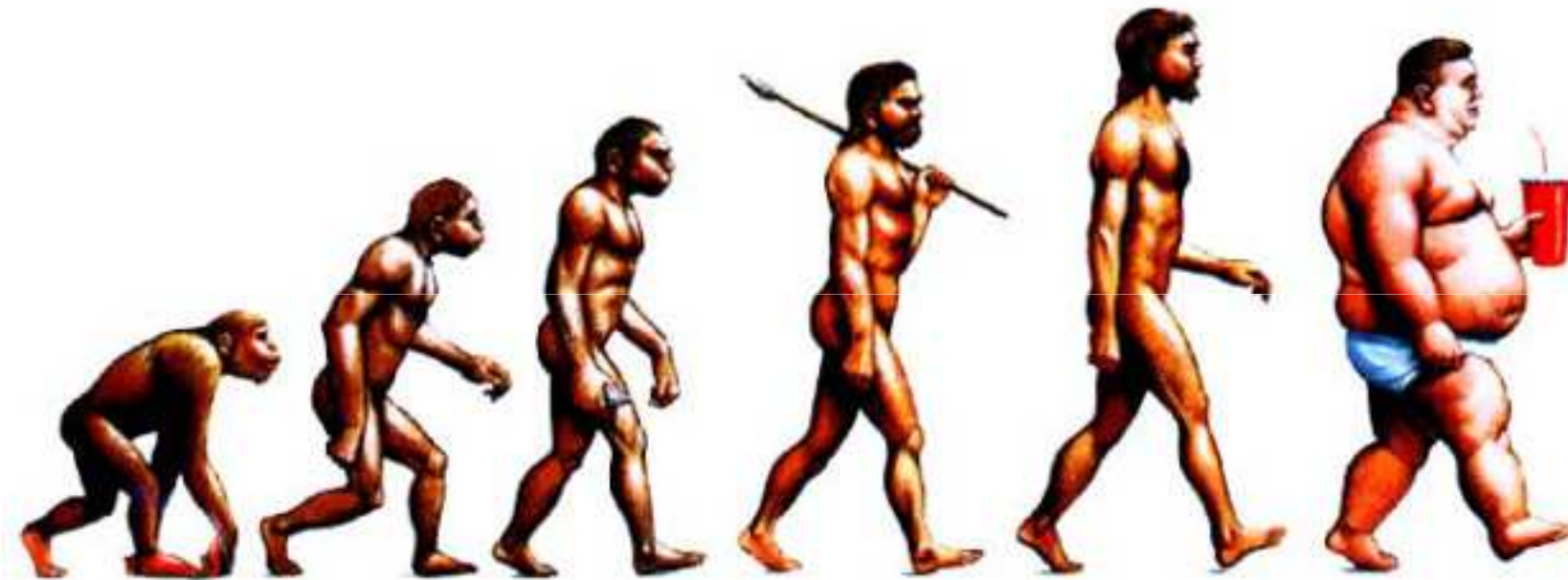
- Don't bet a ridiculous amount!
- As long as we don't have 2000 BTC in this bet, we will simply NOT yet know if bitcoin ECC is broken...

<https://www.betmoose.com/bet/bitcoin-cryptography-broken-in-2015-791>

- Don't expect that code breakers who can make 725,000 \$ elsewhere, will even try to break bitcoin Elliptic Curve
- They would rather steal some bitcoins
  - Possible only if your public key is revealed
  - => Tip: use each Bitcoin address only once!

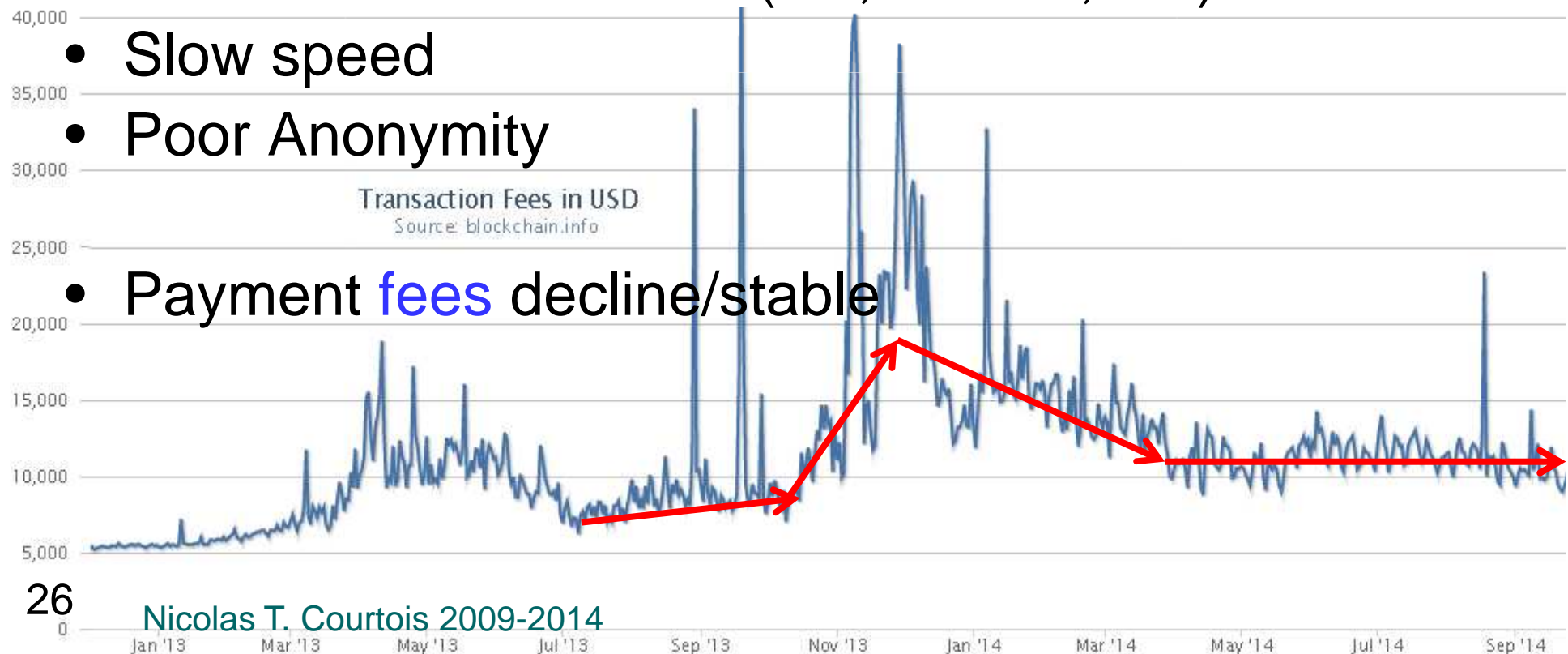


## Is Bitcoin Improving?



## Bitcoin Troubles

- Crypto gets broken?
- Monetary policy: genius, weird or mad?
- 51% attacks and double spending: easy!
- P2P network in decline (XX,000=>5,000)



## So Far...

- Bitcoin has yet failed to achieve the most basic goal: being a decentralized P2P currency (10 major pools control 75%)

## 51% Attacks

See

Nicolas Courtois: [On The Longest Chain Rule and Programmed Self-Destruction of Crypto Currencies](http://arxiv.org/abs/1405.0534) <http://arxiv.org/abs/1405.0534>

**Researcher: cryptocurrencies such as Bitcoin are  
programmed to self destruct**

Posted By: **MrFusion** [[Send E-Mail](#)]

Date: Saturday, 10-May-2014 23:05:41



*Politically Incorrect News  
Stranger than Fiction  
Usually True!*

## Better?

- The “Yahoo of cryptocurrencies” is now waiting for the “Google of cryptocurrencies” to **steal Bitcoin business** purely on technical superiority and without a single hostile shot.

## Better Security Will Prevail?

NOT obvious, and **even**  
**LESS** obvious in financial systems.

A right amount of insecurity:

- allows you to sell insurance,
- trains our survival and cybersecurity skills,
- creates lots of interesting jobs for our students,
- possibly avoids criminals to engage in “more violent” crime...

## Better “Money” Will Prevail?

Crypto engineers like us  
sometimes naively hope that  
“better” currencies will drive  
“not so good” currencies out of business.

In fact the Gresham-Copernicus Law [1517]  
says exactly otherwise!

**Bad currencies** DO frequently drive better  
currencies out of business.

## Better “Money” Will Prevail?

The “bad” option is also happening with bitcoin: it has gained excessive popularity

NOT because it was technically very good (it never was) or had solid intrinsic value, or it was fast and convenient (it never was).

It has thrived because it has created huge expectations which temporarily bitcoin competitors could not meet.

Bitcoin remained the obvious choice, a sort of natural monopoly.



## Network Effects!

Antonopoulos [former UCL student]

points out that

"when you have a technology that is  
'good enough' that achieves network scale [...]  
good enough suddenly becomes perfect"

"I don't see any altcoin displacing it", he says.

If bitcoin crashes, again according to Antonopoulos it will  
be rather because "we blow it up by accident".

[L.A. Bitcoin Meetup Jan 2014]



## Our Works on Bitcoin



-cf. also [blog.bettercrypto.com](http://blog.bettercrypto.com)

- Nicolas Courtois, Marek Grajek, Rahul Naik: [The Unreasonable Fundamental Incertitudes Behind Bitcoin Mining](http://arxiv.org/abs/1310.7935), <http://arxiv.org/abs/1310.7935>
- Nicolas Courtois, Marek Grajek, Rahul Naik: [Optimizing SHA256 in Bitcoin Mining](#), CSS 2014.
- Nicolas Courtois, Lear Bahack: [On Subversive Miner Strategies and Block Withholding Attack in Bitcoin Digital Currency](http://arxiv.org/abs/1402.1718) <http://arxiv.org/abs/1402.1718>
- Nicolas Courtois: [On The Longest Chain Rule and Programmed Self-Destruction of Crypto Currencies](http://arxiv.org/abs/1405.0534) <http://arxiv.org/abs/1405.0534>
- Nicolas T. Courtois, Pinar Emirdag and Daniel A. Nagy: [Could Bitcoin Transactions Be 100x Faster?](#) In proceedings of SECRIPT 2014, 28-30 August 2014, Vienna, Austria.
- Nicolas T. Courtois, Pinar Emirdag and Filippo Valsorda: [Private Key Recovery Combination Attacks: On Extreme Fragility of Popular Bitcoin Key Management, Wallet and Cold Storage Solutions in Presence of Poor RNG Events](http://eprint.iacr.org/2014/848), 16 Oct 2014, <http://eprint.iacr.org/2014/848>
- Poster: [http://www.nicolascourtois.com/bitcoin/POSTER\\_100x\\_Secrypt2014\\_v1.0.pdf](http://www.nicolascourtois.com/bitcoin/POSTER_100x_Secrypt2014_v1.0.pdf)

# Cryptome Renamed My Paper:

# CRYPTOME

Donate for the Cryptome Archive of over 81,300 files from June 1996

key. (Local search temporarily disabled, use Google)

Bitcoin: 1P11b3Xkgagzex3fYusVcJ3ZTVsNwwnrBZ

<http://cryptome.org/2014/05/bitcoin-suicide.pdf> ??????????

=> Actually I show that quite possibly  
bitcoin is EXEMPT from destruction [natural monopoly].

=> Whatever is Bad with bitcoin is  
even worse with most alt-coins.



# Bitcoin vs. Security Engineering

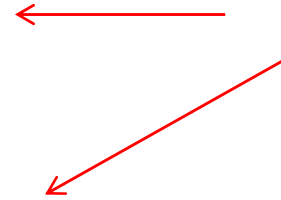


# Re-Engineering Bitcoin:

We postulate:

1. Open design.
2. Least Common Mechanism
3. Assume that attacker controls the Internet  
[Dolev-Yao model, 1983].
4. The specification should be engineered in such a way that it is hard for developers to make it insecure on purpose (e.g. embed backdoors in the system).

**[Saltzer and  
Shroeder 1975]**



## Least Common Mechanism

Violated in Bitcoin also because it uses:

- Open SSL and other standard libraries with massive amounts of code which is not useful at all for bitcoin
- when using TOR
- etc..

# Open Design Principle

[Saltzer and Schroeder 1975]

## Open Design $\neq$ Open Source

Examples: cryptography such as SHA256 (used in bitcoin) is open source but NOT open design – it was designed behind closed doors!





# Open Source vs. Closed Source and Security

## Secrecy:

Very frequently  
an obvious  
business decision.



- Creates entry barriers for competitors.
- But also defends against hackers.

## Kerckhoffs' principle: [1883]

“The system must remain secure should it fall in enemy hands ...”



## Kerckhoffs' principle: [1883]

Most of the time: incorrectly understood.

Utopia. Who can force companies to publish their specs???

No obligation to disclose.

- Security when disclosed.
- Better security when not disclosed.

Yes (1,2,3,4):

1. Military:  
layer the defences.



Yes (2):

2)

Basic economics:

these 3 extra months

(and not more ☹)

are simply worth a  
a lot of money.



Yes (3):

3)

Prevent the erosion of profitability  
/ barriers for entry  
for competitors /  
“inimitability”



Yes (4):

4)

## Avoid Legal Risks

- companies they don't know where their code is coming from, they want to release the code and they can't because it's too risky!
  - re-use of code can COMPROMISE own IP rights and create unknown ROYALTY obligations (!!!)
  - clone/stolen code is more stable, more reliable, easier to understand!





# What's Wrong with Open Source?

## Kerckhoffs principle:

- Rather WRONG in the world of smart cards/HSM...
  - Reasons:
    - side channel attacks,
    - PayTV card sharing attacks
- But could be right elsewhere for many reasons...
  - Example:
    - DES,AES cipher, open-source, never really broken
    - KeeLoq cipher, closed source, broken in minutes...



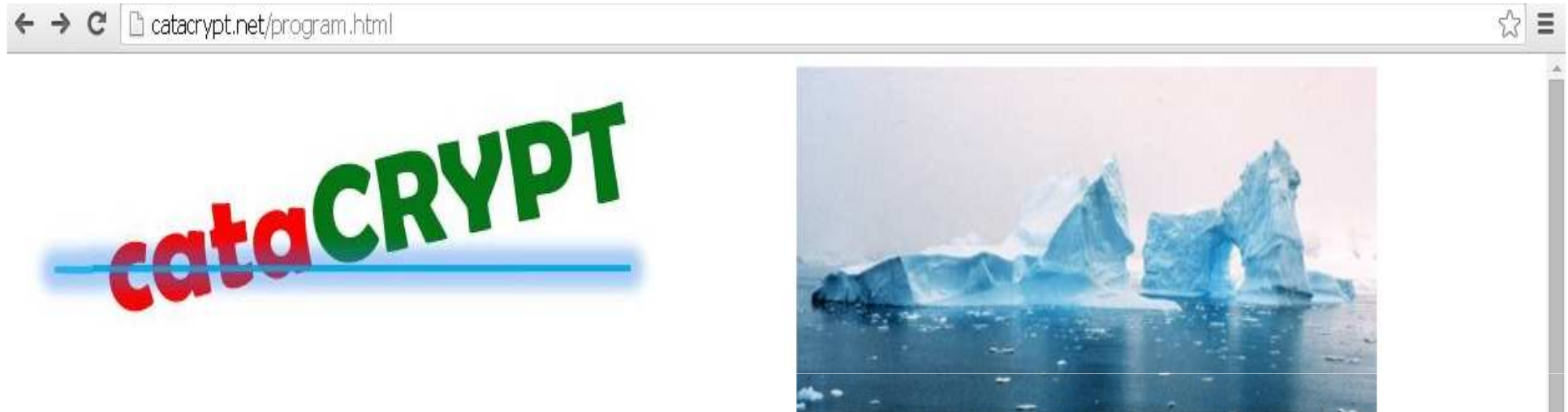
## \*Kerckhoffs principle vs. Public Key Crypto vs. Financial Cryptography

- In Public Key Cryptography one key **CAN** be made public. In practice this means that
  - some **group** of people has it
  - **NO obligation** to disclose, to make it really public (and it is almost never done in serious financial applications)
- Full disclosure for public keys is unbelievably stupid...
  - cf. next slide!

## Do NOT Disclose Public Keys!

- Full disclosure for public keys is simply BAD security engineering and BAD security management.
- Examples:
  - ATMs have like 6 top-level public keys, not really public though
  - in Bitcoin: the public key can remain a secret for years, only a hash is revealed, this is BRILLIANT key management which makes Bitcoin MUCH more secure that it would otherwise be!
  - it does solve the problem raised by Diffie at CataCrypt in San Francisco:  
HOW DO YOU PROTECT AGAINST UNKNWOWN ATTACKS?

# CataCrypt Conference



Workshop on **cata**strophic events related to **crypt**ography and their possible solutions

## Technical Program

[Home](#)

[Committees](#)

[Call for contributions](#)

[Program \(schedule\)](#)


|               |  |
|---------------|--|
|               | <p><b>Venue: Grand Hyatt San Francisco, Union Square, 345 Stockton Street, downtown San Francisco: room Fillmore A - Theatre Level <a href="http://grandsanfrancisco.hyatt.com">http://grandsanfrancisco.hyatt.com</a></b></p> <p><b>October 29, 2014 (together with <a href="#">IEEE Conference on Communications and Network Security (CNS)</a>)</b></p> |
| 08:15 – 08:25 | Opening Remarks: <b>Jean-Jacques Quisquater</b> (UCL, Belgium)   |

# Introducing Bitcoin



## Bitcoin In A Nutshell



- bitcoins are cryptographic tokens, binary data = 010100110101010...
    - stored by people on their PCs or mobile phones
  - ownership is achieved through digital signatures:
    - you have a certain cryptographic key, you have the money.
    - publicly verifiable, only one entity can sign
- 
- An illustration of a hand in a blue sleeve holding a blue pen, signing a yellow document. The document has some black scribbles on it. The entire illustration is surrounded by a pink starburst effect.
- consensus-driven, a distributed system which has no central authority
    - **a major innovation:** financial transactions CAN be executed and policed without trusted authorities.
    - bitcoin is a sort of financial cooperative or a distributed business.
  - based on self-interest:
    - a group of some 100 K people called bitcoin miners own the bitcoin “infrastructure” which has costed > 1 billion dollars (my estimation)
    - they make money from newly created bitcoins and fees
    - at the same time they approve and check the transactions.
    - a distributed electronic notary system

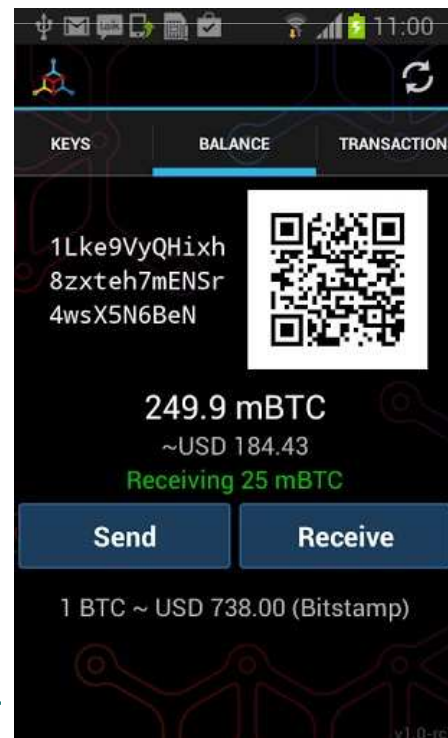


## Two Key Concepts

- initially money are attributed through **Proof Of Work (POW)** to one public key A
  - to earn bitcoins one has to “work” (hashing) and consume energy (pay for electricity)
  - now in order to cheat one needs to work even much more (be more powerful than the whole network), more precisely:
- money transfer from public key A to public key B:
  - **like signing a transfer in front of one notary which confirms the signature,**
  - multiple confirmations: another notary will re-confirm it, then another, etc...
  - we do NOT need to assume that ALL these notaries are honest.
    - at the end it becomes too costly to cheat



## In Practice



## Wallets

- **Wallet**: file which stores your “money”.
- A Bitcoin client App is also called **a wallet**



## Digital Currency

Bitcoin is a

=>PK-based Currency:

- bank account = a pair of public/private ECDSA keys
- spend money = produce a digital signature



## Main Problem:

Bitcoins can be “spent twice”.

Avoiding this “Double Spending” is the main problem when designing a digital currency system.

# Block Chain

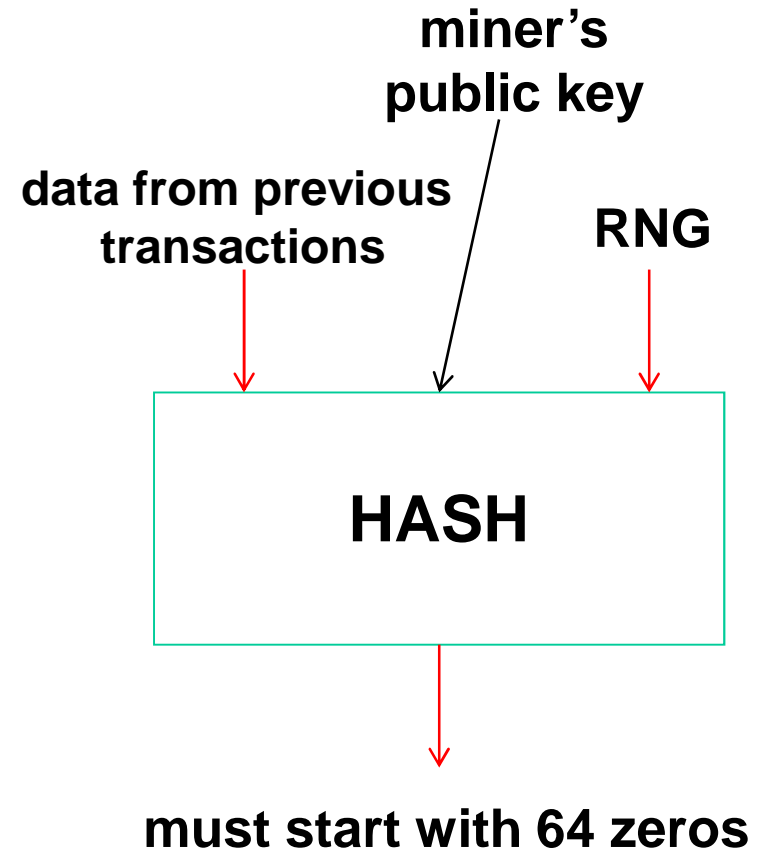


## Bitcoin Mining

- Minting: creation of new currency.
- Confirmation+re-confirmation of older transactions

Ownership:

- “policed by majority of miners”:



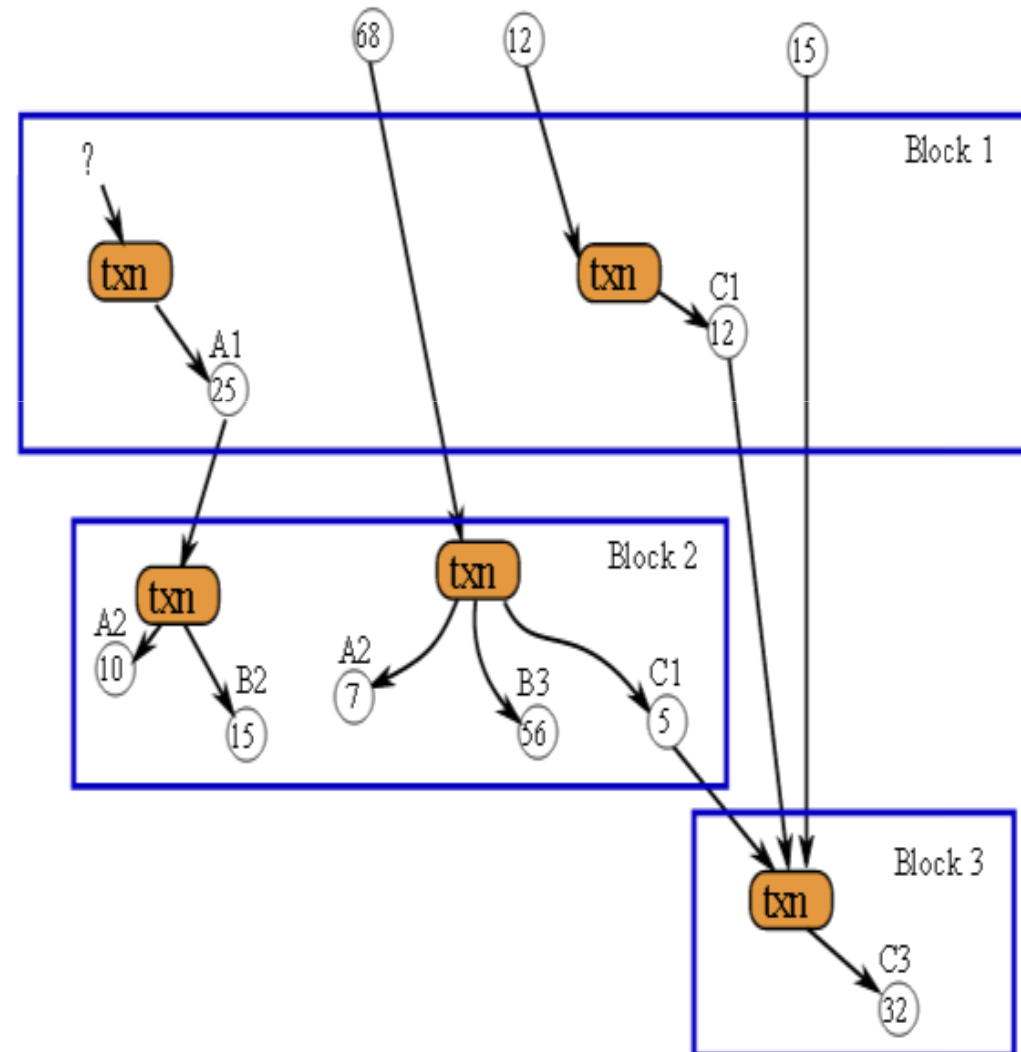
## Block Chain

Def: 

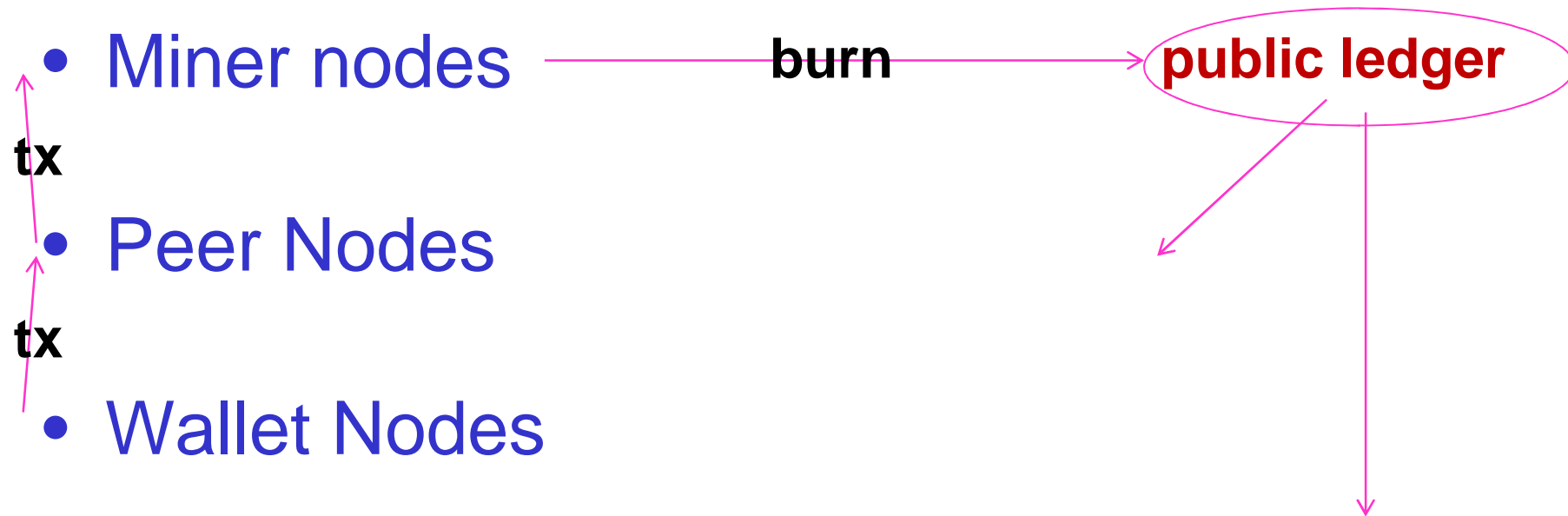
A transaction database  
shared by everyone.

Also a ledger.

Every transaction  
since ever is public.



## Tx LifeCycle





# Bitcoin Address

To: 1K2CcfWYW5sBL2xSeQWXpcmjPCgoXdi36  
Amount: 1.0 BTC

## Ledger-Based Currency

A “Bitcoin Address” = a sort of equivalent of a bank account.

Remarks:

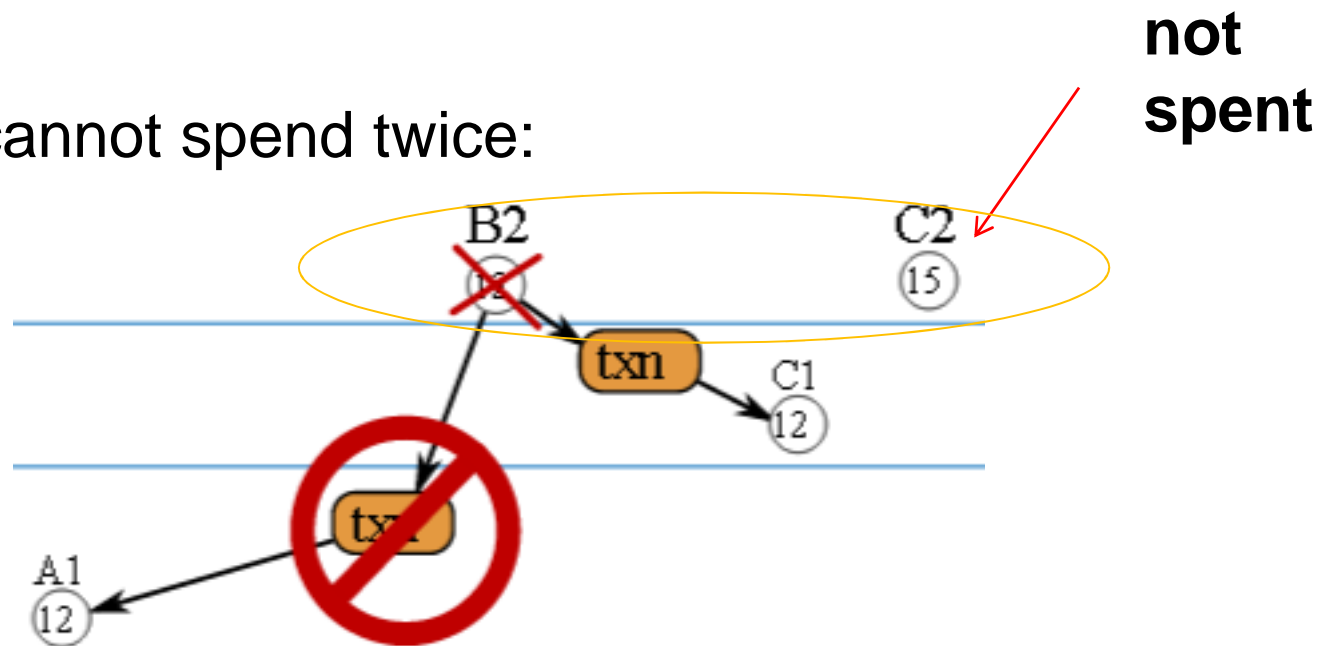
- PK is NOT public!
- only  $H(\text{public key})$  is revealed!
- PK remains confidential until some money in this account is spent.
- SK = private key: always keep private, allows transfer of funds.

## Bitcoin Ownership

Amounts of money are attributed to public keys.

Owner of a certain “Attribution to PK” can at any moment transfer it to some other PK (== another address).

Destructive, cannot spend twice:



# \*Multi-Signature Addresses

## Special Type of Addresses

Bitcoin can require **simultaneously** several private keys,  
in order to transfer the money.

The keys can be stored on different devices (highly secure).

2 out of 3 are also already implemented in bitcoin.

(1 device could be absent, money can still be used).

Very cool, solves the problem of insecure devices...

## Adding Another Layer Of Security

MultiSig:

For example 2 out of 3 signatures are required to spend bitcoins.

# Multi-Sig Concept is NOT new...

**1993**

**Efficient multi-signature schemes  
for cooperating entities**

Olivier Delos <sup>1</sup> and Jean-Jacques Quisquater <sup>2</sup>

# Bitcoin Circulation

To: 1K2CcfWYW5sBL2xSeQWXpcmjPCgoXdi36  
Amount: 1.0 BTC



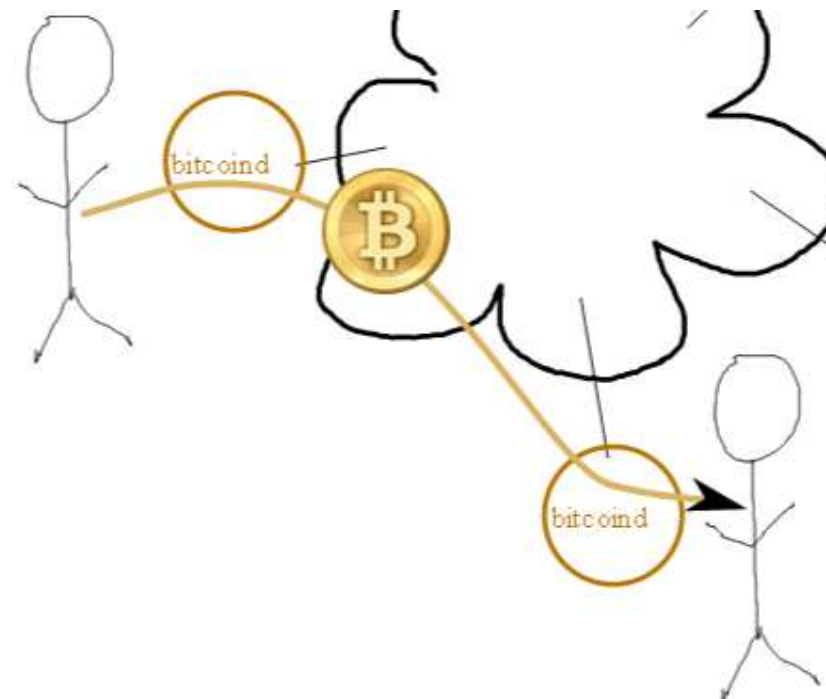
## Bitcoin Transactions:

- between any two addresses [and any two network nodes],
  - at any time [no market closing hours].
  - **validated within 10-60 minutes.**
    - should wait longer for larger transactions, beware of “cheating miners”...
    - 0-confirmation =
      - many websites accept instantly,
      - they trust your application not to double spend
      - and trust miners to reject the second spent based on later time and wider circulation, quite plausible!

## Transfer

To: 1K2CcfWYW5sBL2xSeQWXpcmjPCgoXdi36  
Amount: 1.0 BTC

**SEND**

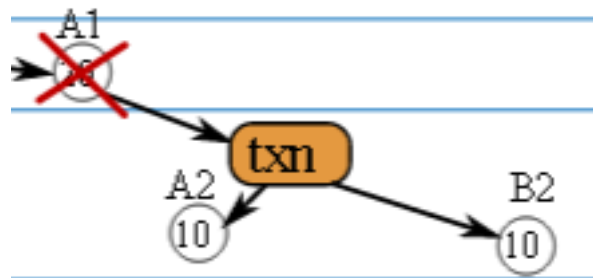


## In / Out

Owner of a certain “Attribution to PK” can at any moment transfer it to some other PK addresses.

=> 0 inputs possible if minting transaction... new money.

=> Several outputs are a norm for bitcoin transactions.



on this picture we  
ignore the fees

## Bitcoin Transfer

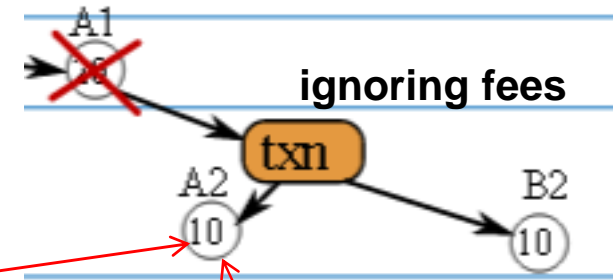
Owner of a certain “Attribution to PK” can at any moment transfer it to any other PK address.



## Attributions

### DEFINITION

“Attribution to PK” =  
act of an owner of  
a previous attribution (always destroyed)  
which transfers a certain amount to the new PK = A2  
(using a digital signature)

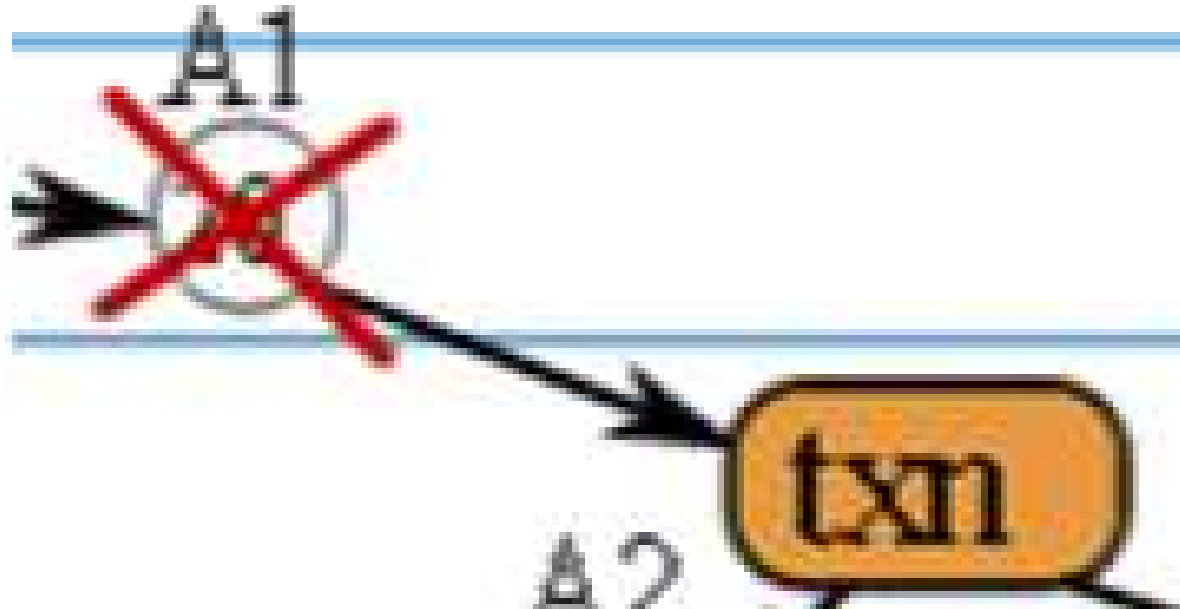


Caveat: Each attribution can be traced back to the initial mining event.

## Fragmentation and Summation Rule

Each PK has a balance, say 20 BTC  
current balance =  $\text{sum}(\text{unspent attributions})$ .

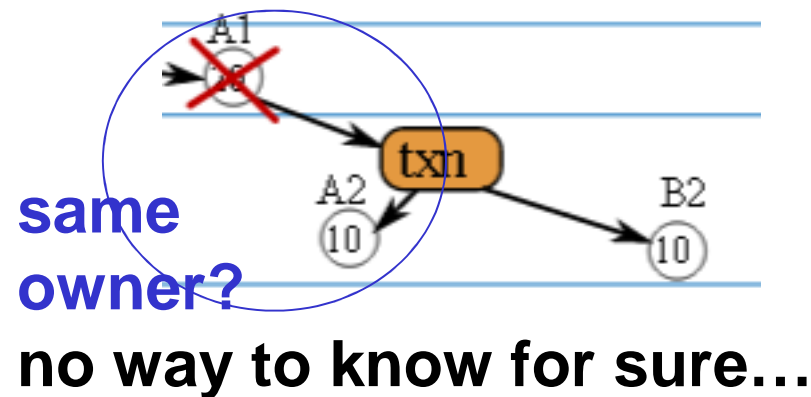
Attributions are ALWAYS destroyed when used,



## From Single Attribution

### Example

- Change: return some money to ourselves inside the same transaction
  - this implies most transactions have 2 or more outputs
  - most apps use the same address
  - could use another fresh address for better anonymity, but too lazy...





## With Multiple Attributions



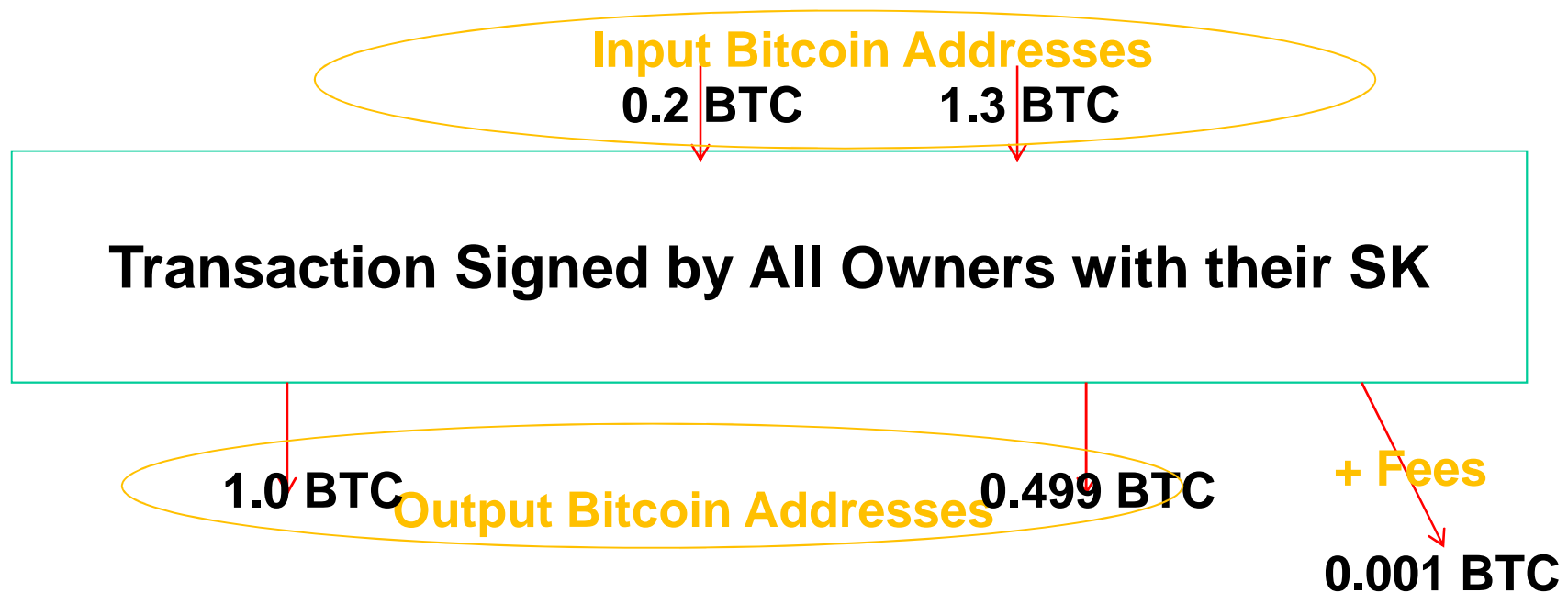
A screenshot of a Bitcoin transaction form. It has a light green background with a black border. The text 'To: 1K2CcfWYW5sBL2xSeQWXpcmjPCgoXdi36' is on the first line, and 'Amount: 1.0 BTC' is on the second line. A blue button with the word 'SEND' in white capital letters is positioned at the bottom right of the form.

To: 1K2CcfWYW5sBL2xSeQWXpcmjPCgoXdi36  
Amount: 1.0 BTC  
SEND

**typical case, even for a single user**

## Bitcoin Transfer

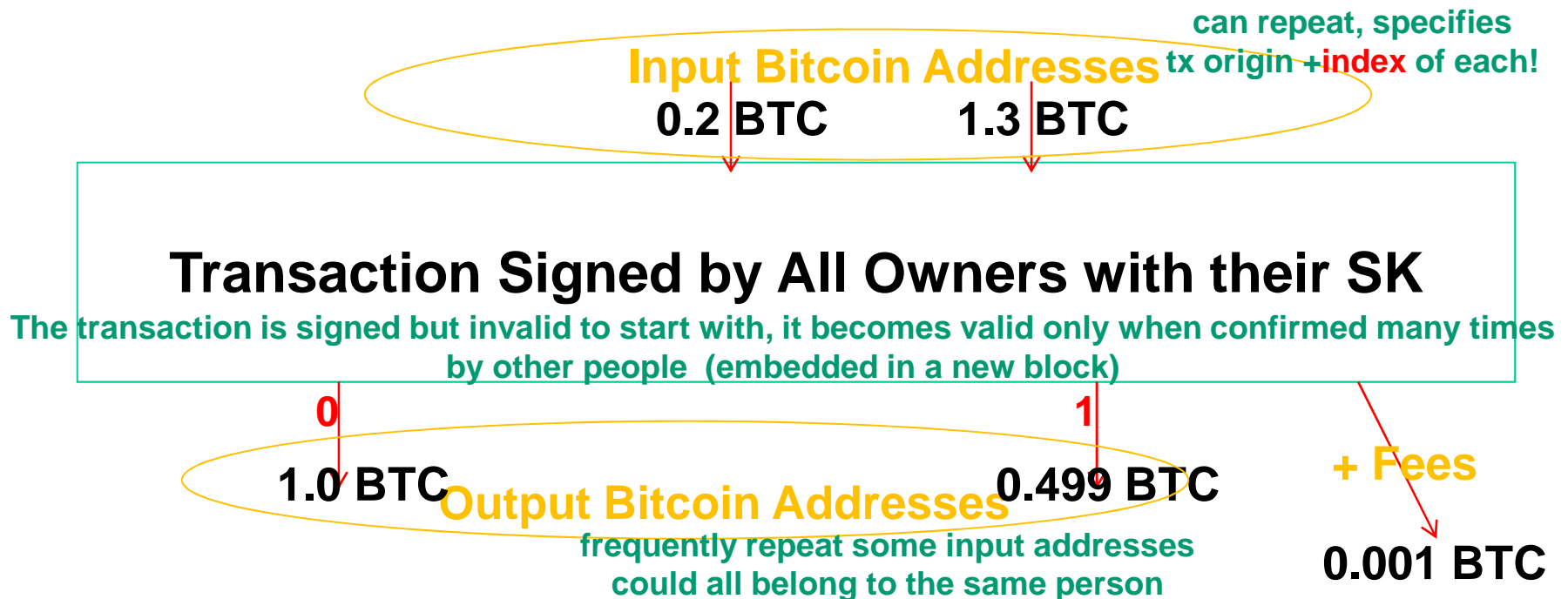
Transactions have multiple inputs and multiple outputs.



## Bitcoin Transfer

Transactions have multiple inputs and multiple outputs.

- helps for anonymity.
- destroys all current attributions,
- requires everybody's signature



## Example 1

### Transaction View information about a bitcoin transaction

99929d9ad149047ae79998592241ddd7ef4ae2f4bb4e057e9c36c4cefa88830

1EWJJCnBuyQDPwVHuCycUCMHCvXTSGLBvk

1MisJY7KwjnhmdaMwyH6v1A3jDQpty7rdg

can repeat,  
tx origin + index of each is  
included in the rawtx



1BaQzo1SyRXZRhQwSvsQJKAUvi5tu3L9uQ

10 mBTC

1rpU1Wa3pYeuJEbRPMWDDCxeH5PDMBrQ9

83.50001 mBTC

1BSy1ARBQfT9PRDYYB6DvzRkbSVRrgbaX3

1.39661 mBTC

can repeat input addresses

94.89662 mBTC

| Summary            |   |
|--------------------|---|
| Size               | 471 (bytes)   |
| Received Time      | 2013-07-20 19:00:32                                     |
| Included In Blocks | <a href="#">247599</a> (2013-07-20 19:03:29 +3 minutes) |
| Confirmations      | 3712 Confirmations                                      |
| Relayed by IP      | <a href="#">5.164.198.173</a> (whois)                   |
| Visualize          | <a href="#">View Tree Chart</a>                         |

| Inputs and Outputs       |   |
|--------------------------|---|
| Total Input              | 95.39662 mBTC                               |
| Total Output             | 94.89662 mBTC                               |
| Fees                     | 0.5 mBTC                                    |
| Estimated BTC Transacted | 94.89662 mBTC                               |
| Scripts                  | <a href="#">Show scripts &amp; coinbase</a> |

## Example 2 = Raw Transaction

```
{
  "hash": "9837485da283ce8ceb0570e2950bb65ebacef9ebd97f871da268d73ea79292c4",
  "ver": 1,
  "vin_sz": 1,
  "vout_sz": 2,
  "lock_time": 0,
  "size": 257,
  "in": [
    {
      "prev_out": {
        "hash": "ba250a395cf37e2d112859ecl4379a605a6fd8e96b406c4f69901abc05d5b47",
        "n": 1
      },
      "scriptSig": "304402206dcf0ef7ca4bfa573ed8f3dc94dca42f5ea46827e8885056d3dfede88e52d49b022077055f3d3c125cc"
    }
  ],
  "out": [
    {
      "value": "5.00000000",
      "scriptPubKey": "OP_DUP OP_HASH160 ddc1120deb91acda0d3e5774a2b8908e3424f532 OP_EQUALVERIFY OP_CHECKSIG"
    },
    {
      "value": "13.07598401",
      "scriptPubKey": "OP_DUP OP_HASH160 88f1271342d5f2202995c6e74ed07b81caec7633 OP_EQUALVERIFY OP_CHECKSIG"
    }
  ]
}
```

unique ID on 256 bits = the hash of the whole

list of input attributions:  
origin tx, index n, ECDSA signature

list of output attributions

amount BTC

H(recipient PK)

## Remarks:

About XX million transactions ever made.

To know the balance of one account, we must “in theory” store ALL the transactions which send money for this address and then check ALL transactions made since then to see some of these are not already spent.

Full bitcoin network nodes stored all transactions ever made and checks their correctness (all the digital signatures).

About 24 Gbytes data, 48 hours typical download.

In practice one could skip check for things confirmed by many miners... dangerous though. There is no absolute proof that miners have already checked them (maybe they forgot, a bug).

# Transaction Scripts

## \*\*\*Scripts

```
{
  "hash": "9837485da283ce8ceb0570e2950bb65ebacef9ebd97f871da268d73ea79292c4",
  "ver": 1,
  "vin_sz": 1,
  "vout_sz": 2,
  "lock_time": 0,
  "size": 257,
  "in": [
    {
      "prev_out": {
        "hash": "ba250a395cf37e2d112859ecl1d4379a605a6fd8e96b406c4f69901abc05d5b47",
        "n": 1
      },
      "scriptSig": "304402206dcf0ef7ca4bfa573ed8f3dc94dca42f5ea46827e8885056d3dfede88e52d49b022077055f3d3c125cc"
    }
  ],
  "out": [
    {
      "value": "5.00000000",
      "scriptPubKey": "OP_DUP OP_HASH160 ddc1120deb91acda0d3e5774a2b8908e3424f532 OP_EQUALVERIFY OP_CHECKSIG"
    },
    {
      "value": "13.07598401",
      "scriptPubKey": "OP_DUP OP_HASH160 88f1271342d5f2202995c6e74ed07b81caec7633 OP_EQUALVERIFY OP_CHECKSIG"
    }
  ]
}
```

**Signature Script**

list of output attributions

**Redemption Script**

**0** → "value": "5.00000000",  
"scriptPubKey": "OP\_DUP OP\_HASH160 ddc1120deb91acda0d3e5774a2b8908e3424f532 OP\_EQUALVERIFY OP\_CHECKSIG"

**1** → "value": "13.07598401",  
"scriptPubKey": "OP\_DUP OP\_HASH160 88f1271342d5f2202995c6e74ed07b81caec7633 OP\_EQUALVERIFY OP\_CHECKSIG"

**H(recipient PK)**



# Spot On Signatures

## Signed Tx / Final Tx

byte by byte (similar but not identical to raw blocks seen before)  
(this is done twice, with different scriptSig)

|                 |                                 |   |
|-----------------|---------------------------------|---|
| version         |                                 | 01 00 00 00   |
| input count     |                                 | 01  |
| input           | previous output hash (reversed) | 48 4d 40 d4 5b 9e a0 d6 52 fc a8 25 8a b7 ca a4 25 41 eb 52 97 58 57 f9 6f b5 0c d7 32 c8 b4 81 |
|                 | previous output index           | 00 00 00 00   |
|                 | script length                   | scriptSig length 1 byte, e.g. 25=0x19 or 138=0x8A   |
|                 | scriptSig                       | script containing signature scriptSig   |
|                 | sequence                        | ff ff ff ff   |
| output count    |                                 | 01  |
| output          | value                           | 62 64 01 00 00 00 00 00 (in Satoshis)   |
|                 | script length                   | scriptPubKey length 1 byte, e.g. 25=0x19  |
|                 | scriptPubKey                    | script containing destination address scriptPubKey  |
| block lock time |                                 | 00 00 00 00 (never used so far)   |

2 scripts will be detailed later

$$\text{len}(1i/1o) = 223 = 4+1+32+4+1+ 1+71+ 1+65+ 4+1+8+ 1+25+4$$

## First scriptSig

It is scriptPubKey BUT copied from the previous transaction  
(peculiarity)

**len=** 25=3+20+2 typically

## Second scriptSig

sign+PKey

**len= 1+71+ 1+65 = 138 BUT NOT ALWAYS!**

**scriptSig**

|                    |            |   |  |  |
|--------------------|------------|---|--|--|
| PUSHDATA 47        |            | 47  |  |  |
| signature<br>(DER) | sequence   | 30  |  |  |
|                    | length     | 44  |  |  |
|                    | integer    | 02  |  |  |
|                    | length     | 20  |  |  |
|                    | X <b>r</b> | 2c b2 65 bf 10 70 7b f4 93 46 c3 51 5d d3 d1 6f c4 54 61 8c 58 ec 0a 0f |  |  |
|                    | integer    | 02  |  |  |
|                    | length     | 20  |  |  |
|                    | Y <b>s</b> | 6c 66 24 d7 62 a1 fc ef 46 18 28 4e ad 8f 08 67 8a c0 5b 13 c8 42 35 f1 |  |  |
| SIGHASH_ALL        |            | 01  |  |  |
| PUSHDATA 41        |            | 41  |  |  |
| public key         | type       | 04  |  |  |
|                    | X          | 14 e3 01 b2 32 8f 17 44 2c 0b 83 10 d7 87 bf 3d 8a 40 4c fb d0 70 4f 13 |  |  |
|                    | Y          | 10 f9 81 92 6e 53 a6 e8 c3 9b d7 d3 fe fd 57 6c 54 3c ce 49 3c ba c0 63 |  |  |

**scriptSig1**

**scriptSig2**

## Is Bitcoin Secure?

Satoshi claimed it is...



# Incidents at Operation: Bad Randoms

## Bad Randoms

First publicized by Nils Schneider:

28 January 2013

D47CE4C025C35EC440BC81D99834A624875161A26BF56EF  
7FDC0F5D52F843AD1

⇒ repeated more than 50 times...

Used twice by the SAME user!



## ECDSA Signatures

Let **d** be a private key, integer **mod n** = ECC [sub-]group order.

- Pick a random non-zero integer  $0 < a < n-1$ .
- Compute  $R = a \cdot P$ , where **P** is the base point (generator).
- Let  $r = (a \cdot P)_x$  be its x coordinate.
- Let  $s = (H(m) + d \cdot r) / a \bmod n$ .

The signature of **m** is the pair **(r,s)**.

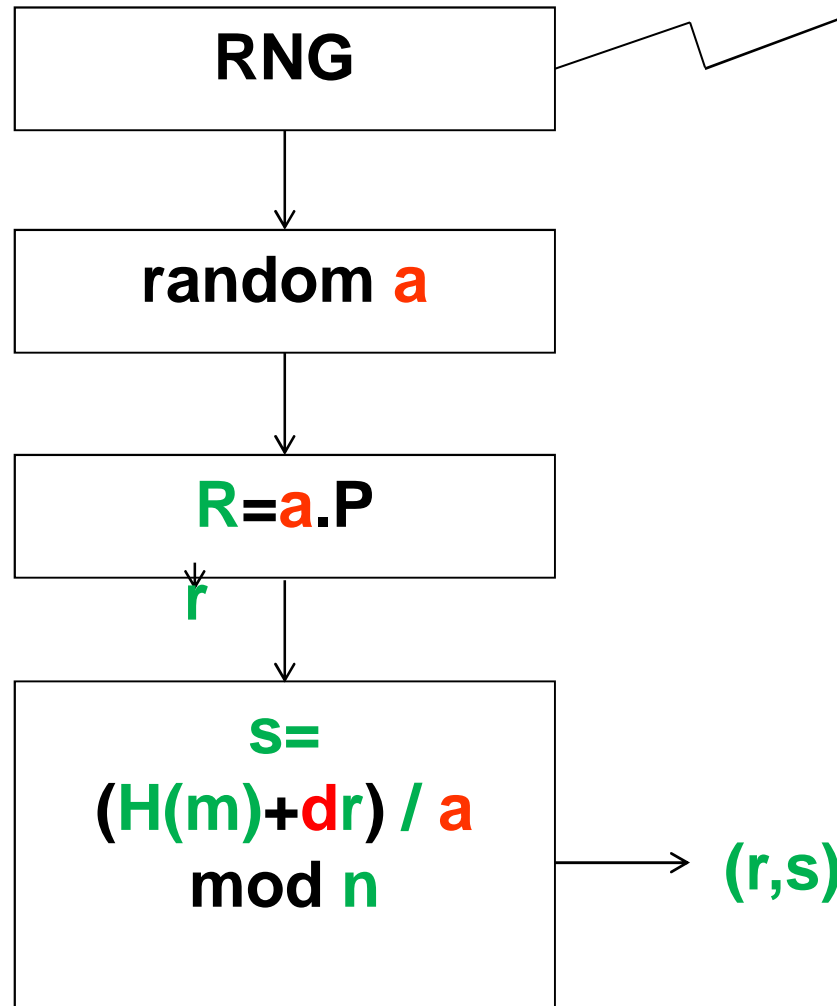
(512 bits in bitcoin)



# Attack – 2 Users

has already happened  
100 times in Bitcoin

random **a**: must be kept secret!



same **a** used twice  $\Rightarrow$   
detected in public  
blockchain  $\Rightarrow$

$$(s_1 a - H(m_1)) / d_1 = r = (s_2 a - H(m_2)) / d_2 \bmod n$$

$\Rightarrow$

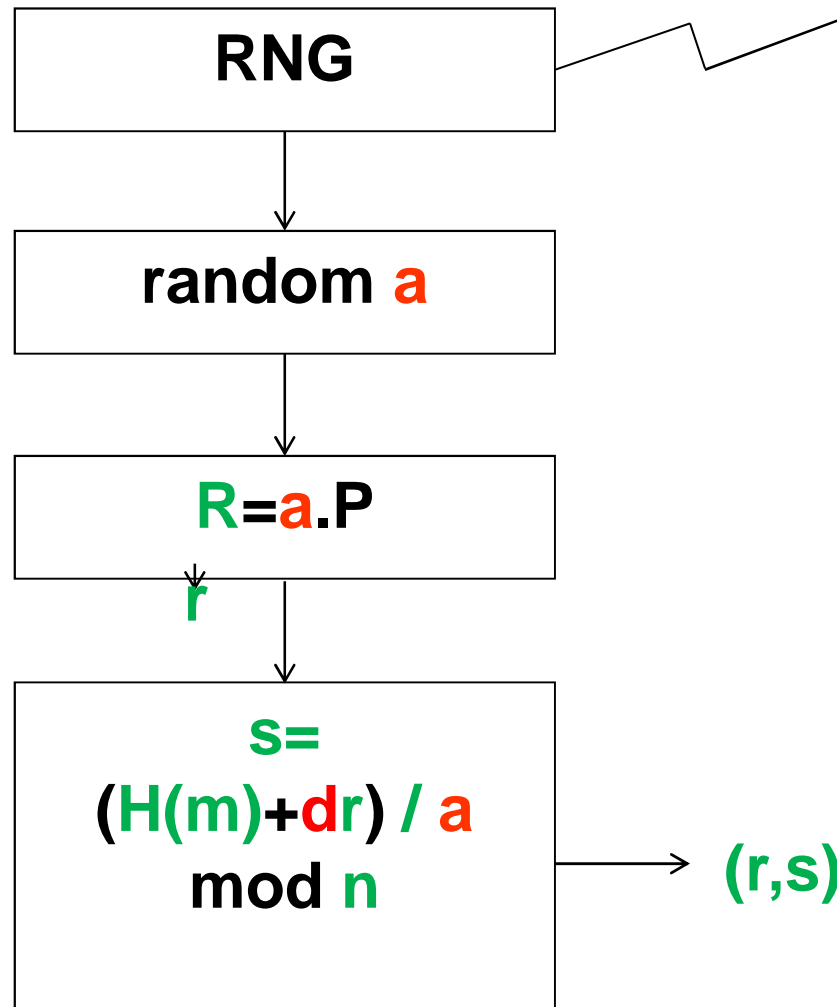
$$r(d_1 - d_2) + a(s_1 - s_2) = H(m_2) - H(m_1) \bmod n$$

each person can steal the  
other person's bitcoins!

$\Rightarrow$  any of them CAN  
recompute **k** used

# Attack – Same User

random **a**: must be kept secret!



has also happened  
100 times in Bitcoin

same **a** used twice by the  
same user ( $d_1 = d_2$ ). In this  
case we have:  $(s_1 a - H(m_1)) =$   
 $rd =$   
 $(s_2 a - H(m_2)) \bmod n$   
 $\Rightarrow a = (H(m_1) - H(m_2)) / (s_1 - s_2)$   
 $\bmod n$  AND now  
 $d = (sa - H(m)) / r \bmod n$

anybody can steal  
the bitcoins!

# Stopped in August 2013

Android bug was fixed...

## Dec. 2013

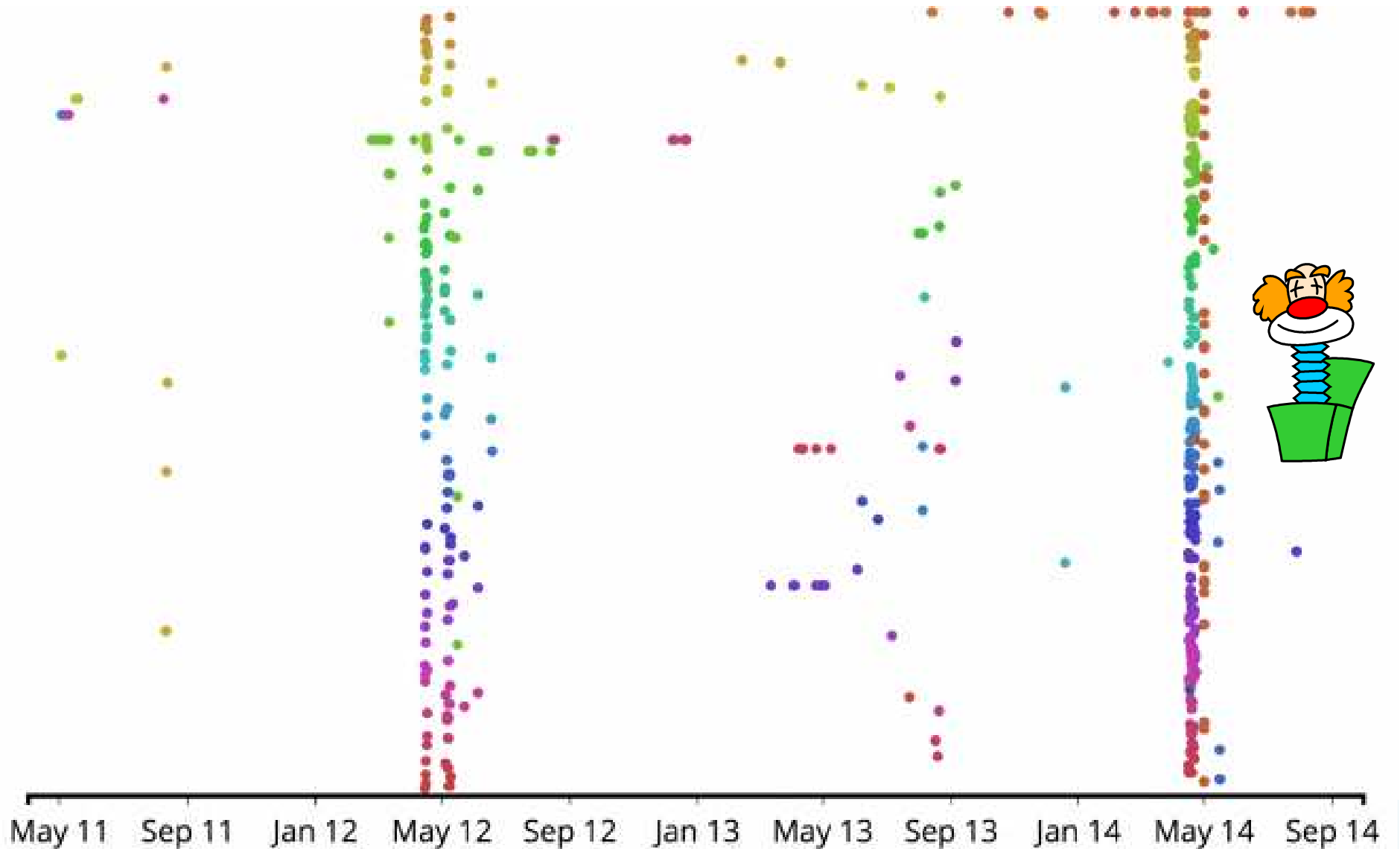
At 30C3 conference in Germany on 28 Dec 2013  
Nadia Heninger have reported that they have identified  
a bitcoin user on the blockchain  
which has stolen some 59 BTC due to  
these bad randomness events,

The money from the thefts is stored at:

<https://blockchain.info/address/1HKywxIL4JziqXrzLKhmb6a74ma6kxbSDj>

Still sitting there, he is NOT trying to spend it...  
too famous? Afraid to be traced and caught?

## Second Major Outbreak – May 2014



## Recent Bad Randoms

From my own scan:

0f25a7cc9e76ef38c0feadcfa5550c173d845ce36e16bde09829a  
3af57097240.

Appears 8 times in block 322925  
28 September 2014

Used by different users...



## So What?

### Previous attacks:

- Classical bad random attacks typically concern only very few bitcoin accounts, and only some very lucky holders of bitcoins can actually steal other people's bitcoins.
- Only **a few hundred accounts** in the whole history of bitcoin are affected.



# The Really Scary Attacks

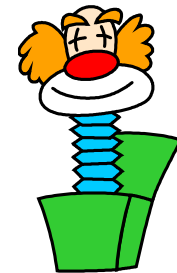
New attacks [Courtois et al. October 2014]

=> under certain conditons

ALL bitcoins in cold storage

can be stolen

=> millions of accounts potentially affected.







## New Paper:

### Private Key Recovery Combination Attacks: On Extreme Fragility of Popular Bitcoin Key Management, Wallet and Cold Storage Solutions in Presence of Poor RNG Events

cf.

[eprint.iacr.org/  
2014/848/](http://eprint.iacr.org/2014/848/)

Nicolas T. Courtois<sup>1</sup>

Pinar Emirdag<sup>2</sup>

Filippo Valsorda<sup>3</sup>

<sup>1</sup> University College London, UK

<sup>2</sup> Independent market structure professional, London, UK

<sup>3</sup> CloudFlare, London, UK



**Abstract.** In this paper we study the question of key management and practical operational security in bitcoin digital currency storage systems. We study the security two most used bitcoin HD Wallet key management solutions (e.g. in BIP032 and in earlier systems). These systems have extensive audit capabilities but this property comes at a very high price. They are excessively fragile. One small security incident in a remote corner of the system and everything collapses, all private keys can be recovered and ALL bitcoins within the remit of the system can be stolen. Privilege escalation attacks on HD Wallet solutions are not new. In this paper we take it much further. We propose new more advanced **combination attacks** in which the security of keys hold in cold storage can be compromised without executing any software exploit on the cold system, but through security incidents at operation such as **bad random number or related random events**.

In our new attacks all bitcoins over whole large security domains can be stolen by people who have the auditor keys which are typically stored in hot systems connected to the Internet and can be stolen easily. Our combination attacks allow to recover private keys which none of the

## Is There a Fix?

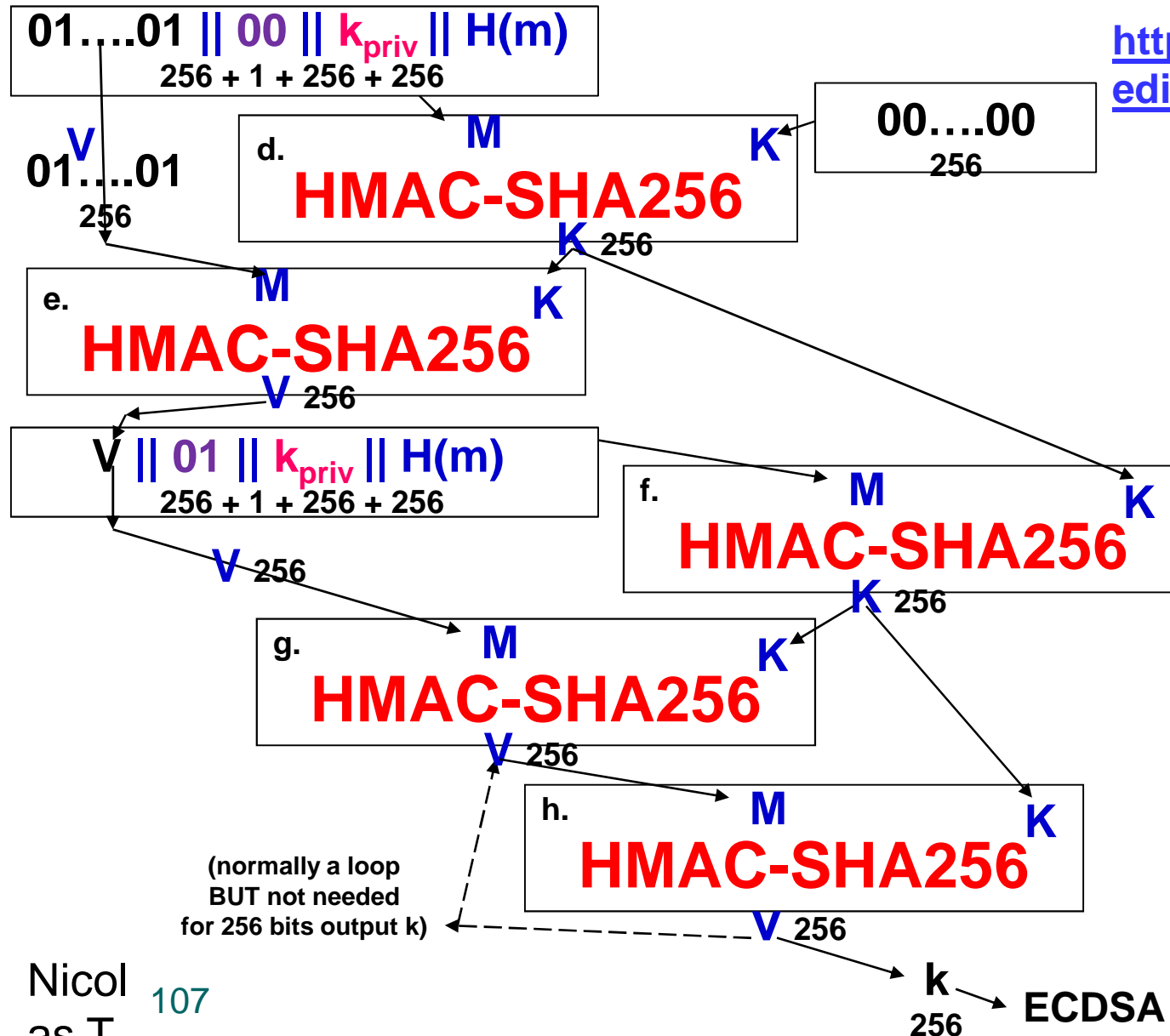
Solution: RFC6979 [Thomas Pornin]

HOWEVER,

no existing cold storage solution  
which have NOT already applied RFC6979  
can claim to resist our attacks.

## RFC6979 [Pornin] = 5+ applications of HMAC

<http://www.rfc-editor.org/rfc/rfc6979.txt>



# Which Systems Are Affected?

Solution: RFC6979 [Pornin]

- Already applied by
  - Electrum, Multibit, Trezor
- Yet unpatched:
  - blockchain.info – insecure,
  - BitcoinD Core – waiting for a patch to be applied,



Details:

a talk at Hack in The Box conference 15/10/2014:

<http://conference.hitb.org/hitbsecconf2014kul/materials/D1T1%20-%20Filippo%20Valsorda%20-%20Exploiting%20ECDSA%20Failures%20in%20the%20Bitcoin%20Blockchain.pdf>