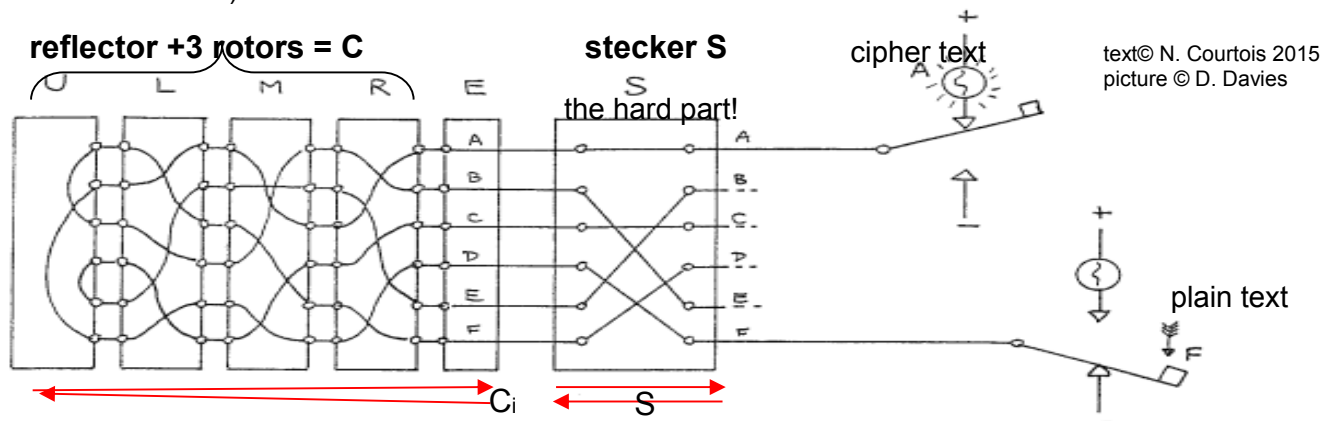


## Understanding HOW Enigma Was Broken or How Mathematicians Won the WW2 [cf. UCL COMPGA18 course].

- In 1920s a “commercial Enigma” machine was commercialized. Rotors rotate and each letter is encrypted by a different circuit of permuted wires. This could be broken by paper and pencil. Since 1930 Germany encrypted their communications with a “military Enigma” which is a LOT more complex. It differs by the introduction of the so called ‘**stecker**’ [plug board] **S** which increased the entropy of the key tremendously by some 47 extra bits (150 million million times) to about 76 bits total.



- For the next 10 years it was considered unbreakable and its details remained unknown at BP. To know the [secret] connections of one rotor was equivalent to having an additional secret of 88 bits. However the same rotors remained in use for decades. The French had a spy H. in the German security service. This spy has not supplied an Enigma machine however he gave them a manual which contained examples settings+plaintext+ciphertext.
- This allowed Polish mathematicians to recover the wirings of the rotors [by mathematics/factoring permutations]. In 1932-1938 majority of messages were decrypted by Poles [very weak Enigma initialization procedures].
- In Dec 1938 the setup method was changed and the Poles invented 2 new attacks: **Rejewski's bombe** which and **Zygalski sheets** method. All the knowledge [and machines] were given to allies [France and Britain] in July 1939.
- Zygalski** method was already very expensive and actually it was only fully implemented by Jeffreys at Bletchley Park a year later, and allowed the Poles+French+UK to decrypt Enigma traffic until May 1940 [invasion of France].
- In late 1939 **Alan Turing** and other started developing a new [+general ++powerful ++expensive] attack which would decrypt Enigma messages ignoring weak setup, anticipating that one day Polish attacks will stop working.
- The cost of breaking Enigma was huge, say a price of 200 extra war planes and **Churchill** personally made sure that code breakers obtained money required. First Bombe by Turing [just before May 1940] didn't work well.
- Another mathematician **Gordon Welchman** has invented a so called diagonal board, a tremendous improvement.
- Overall we call it the **Turing–Welchman Bombe Attack** on Enigma. It allows to totally **ELIMINATE the stecker** and find the 3 rotor positions [previous attacks assumed weak setup/init or less connections in the stecker].

The attack works in 3 stages [it will be explained during our visit and we will see a working real-life demo in block B]:

- First code breakers need to guess a **crib**, part of the plaintext, e.g. OBERKOMMANDODERWEHRMACHT.
- They compared plaintext to many encrypted messages. Enigma has one terrible weakness: no letter would encrypt to itself. Consequently, most of the time it was possible to reject a given alignment of the crib.
- Codebreakers could therefore guess a plausible alignment which had some chances to be the correct one.
- From this they created a so called **menu**. A graph in which pairs of letters are connected, say G encrypts to P.
- Turing approach exploits short cycles in this graph. For example imagine we have a simple cycle of length 3  $G \Rightarrow P \Rightarrow A \Rightarrow G$  [cycles could be of any length, shorter cycles e.g. length 2 were broken by earlier attacks].
- Turing focused on the possibility to obtain the same G again. If we connect several Enigmas in a closed loop [just following our cycle] and if we connect wire G to a battery, the current will come out at same letter G.
- A cycle in the 'menu' implies a **fixed point** for a certain sequence of full encryptions. For example  $P(G)=G$  where **P** is a combination of full Enigmas with settings which differ by a few steps in time, e.g.  $t=t_0+2$  in first Enigma.
- Let **S** be the permutation of the stecker and let **C<sub>2</sub>**, **C<sub>5</sub>** and **C<sub>8</sub>** are the combined permutations of going through 3 rotors and back in Enigma at clocks 2, 5 and 8. We have  $P = S^{-1} \cdot C_2 \cdot S \cdot S^{-1} \cdot C_5 \cdot S \cdot S^{-1} \cdot C_8 \cdot S$  [read right to left].
- So  $G = S^{-1} \cdot C_2 \cdot S \cdot S^{-1} \cdot C_5 \cdot S \cdot S^{-1} \cdot C_8 \cdot S \cdot G$  which simplifies to  $S \cdot G = C_2 \cdot C_5 \cdot C_8 \cdot S \cdot G$ ; **S(G)** is a fixed point for **C<sub>2</sub>·C<sub>5</sub>·C<sub>8</sub>** !
- This permutation **C<sub>2</sub>·C<sub>5</sub>·C<sub>8</sub>** does NOT depend on the stecker and is implemented by a serial connection of 3 so called Letchworth Enigmas [with separated inputs and outputs], each has 3 drums, connections are done at the back of the bombe through 26-wire red cables 'spaghetti'. Welchman diagonal board=>extra connections.
- A nice trick is just to input some random letter say A to this circuit **C<sub>2</sub>·C<sub>5</sub>·C<sub>8</sub>** connected to form an 'infinite' loop.
- Most of the time the machine actually tries totally wrong settings. Then A is not a fixed point for **C<sub>2</sub>·C<sub>5</sub>·C<sub>8</sub>**.
- So current comes out at another letter, then because we have a closed loop, it goes inside again. Typically all 26 letters are 'live'. This is clearly a wrong setting. Frequently we have no fixed point at all, 26 wires are active => all 'fast' rotors in the machine turn to try the next setting [search actually done backwards  $t_0=ZZZ..AAA$ ].
- Now if **C<sub>2</sub>·C<sub>5</sub>·C<sub>8</sub>** has a fixed point D, the current will reach 25 wires EXCEPT D. If this happens the machine stops. We should now think that **S(G)=D**. This is if the 3 rotors are at correct positions. Then recover two full S on the checking machine. Typically it will have 10 swaps and 6 fixed points. This is tried endless times until correct!
- Overall during WW2 Britain have cracked 25,000 keys and decrypted 2.5M messages shortening war by 2 years.